Averting multi-qubit burst errors in surface code magic state factories

Tough Errors

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Background: surface code & magic state distillation

- Surface code: rectangular patch of physical qubits; multiqubit operations by merging and splitting patches
- Can fault-tolerantly perform Clifford operations, but need *T* gate to complete universal gate set
- *T* gate can be performed using a *magic state* $|T\rangle$
- Magic states of high fidelity are generated in *magic state factories* made of surface code tiles
- Magic state factories repeatedly perform *magic state distillation*
 - In this work, we focus on 15-to-1 distillation, which suppresses errors from order p to order p^3

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 Magic state distillation estimated to be 60-95% of total program cost (qubitcycles)



Background: multi-qubit burst errors

- Quantum error correction relies on sufficiently-small and unchanging physical error rates
- Physical error rates fluctuate significantly on current hardware in a variety of ways
- We focus on *multi-qubit burst errors*: many qubits experiencing an increase in error rate at the same time
- Common source of burst errors in superconducting hardware: cosmic ray impacts
 - Rate of rays can be reduced by shielding, but a single burst error could ruin an hours-long computation
 - Gap engineering can reduce direct sensitivity to radiation, but may come with fabrication tradeoffs and does not solve the whole problem

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M. McEwen et al., "Resolving catastrophic error bursts from cosmic rays in large arrays of superconducting qubits," *Nature Physics* (2022)

Noise model: Direct



Based on M. McEwen et al., "Resolving catastrophic error bursts from cosmic rays in large arrays of superconducting qubits," Nature Physics (2022).

- Model: ray reduces T_1 times in some radius r_{CRE} . Qubit error rates increase linearly towards center, with maximum reduction at the center of $f_{T_1}T_1^{init}$.
- Ray impacts are Poisson-distributed with rate Γ. Goes away after some time.

Noise model: TLS Scrambling



Based on C. D. Wilen et al., "Correlated charge noise and relaxation errors in superconducting qubits," Nature (2021) and T. Thorbeck et al., "Two-Level-System Dynamics in a Superconducting Qubit Due to Background Ionizing Radiation," PRX Quantum (2023).

- Model: ray scrambles T_1 times randomly within radius r_{CRE} .
- Ray impacts are Poisson-distributed with rate Γ. Requires active re-calibration to fix.

Baseline: Code expansion

Y. Suzuki et al., "Q3DE: A fault-tolerant quantum computer architecture for multi-bit burst errors by cosmic rays," MICRO 2022

- Allocate extra buffer space around each patch
 - With enough buffer space, can perform distillation in 5d_m steps instead of 6d_m
- Upon burst event, expand patch to increase error resilience



Baseline: Code expansion

- How much buffer space do we need? Depends on cosmic ray parameters
- For **Direct** model, d_{extra} depends on r_{CRE} and f_{T_1}
- For **Scrambling** model, we assume that added distance must be sufficient for worst-case set of broken qubits, so $d_{\text{extra}} = 2r_{CRE}$
- Assume *d*_{extra} must be doubled if there is a significant chance of two simultaneous events

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• Depends on $\Gamma \times T_{\text{offline}}$



Baseline: Distributed

Q. Xu et al., "Distributed Quantum Error Correction for Chip-Level Catastrophic Errors," Phys. Rev. Letters (2022)

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- Encode each logical qubit in higher-level distributed code
- A detected burst error is treated as a heralded erasure error (assume entire patch is broken)
- A code with distance d qubits can tolerate d 1 simultaneous erasures
 - $\Gamma \times T_{\text{offline}}$ (probability of simultaneous events) determines required higher-level code



Solution: partially-offline magic state factories

- Magic state factories do not store long-term logical information; we do not have to protect them as carefully
- Idea: if a ray hits, just turn parts of the factory offline until recovery
- Re-mapping allows factory to operate even under more severe disruption

Normal operation



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Mitigating burst errors in magic state factories



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Comparison to baselines

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Both baselines assume **instant** and **complete** detection of burst events, so we compare ۲ under that assumption



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Realistic detection of burst error events

- How quickly can we reliably detect burst errors when our only information is QEC error syndromes?
- Count error syndromes in spatiotemporal windows
- Define spatial windows of size $w_s \times w_s$
- For each spatial window, determine average baseline syndrome rate per stabilizer $p_{\text{syn},i}$
- Define temporal window size w_t and set a threshold number of counts $n_{\text{th},i}$ based on desired false positive rate (FPR)
- Each cycle, count syndromes in each window. If the count exceeds $n_{\text{th},i}$, a detection event is triggered
- Upon a detection event, turn off all qubits with radius $r_{\rm off}$ of the window for duration $T_{\rm offline}$



Burst error detection latency



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Overhead of re-mapping under realistic detection

- Fixed temporal overhead: $\Gamma \times T_{\text{offline}} = 10^{-5}$
- Spatial overhead determined by *T* buffer size, which is set by reliable-detection latency
- **Direct** model: latency determined by f_{T_1} and r_{CRE}
 - Less than 2x overhead for most of the studied parameter space, but quickly grows for small and weak rays
- Scrambling model: latency determined by *r*_{CRE}

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 Reliable detection is difficult; need to design for worst-case ray

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• Overhead quickly grows as r_{CRE} decreases



Summary and discussion

- By tailoring burst error mitigation to magic state factories, we reduced mitigation overheads by 6.5-13.9× compared to previous methods
- Scales favorably with Γ (no time overhead until an event happens)
- Easily extends to different magic state factory layouts
 - Overhead factor will *decrease* with increasing factory size, while baseline overheads will *increase*
- Re-mapping factories may be useful for other error sources (fluctuating TLSs, calibration drift, etc.)
- Our method does not apply to logical program qubits still need a larger-overhead mitigation method for some parts of the processor
- Detection of weaker burst errors is more difficult than previously assumed we need to carefully study implications for compute qubits